

# The density of truth in monadic fragments of some intermediate logics

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## Abstract

This paper is an attempt to count the proportion of tautologies of some intermediate logics among all formulas. Our interest concentrates especially on Dummett's and Medvedev's logics and their  $\{\rightarrow, \vee, \neg\}$  fragments over language with one propositional variable.

## 1 Introduction

Let  $L$  be some logical calculus. Let  $|T_n|$  be a number of tautologies of length  $n$  of that calculus and  $|F_n|$  be a number of all formulas of that length. We define the density  $\mu(L)$  as:

$$\mu(L) = \lim_{n \rightarrow \infty} \frac{|T_n|}{|F_n|}$$

The number  $\mu(L)$  if exists, is an asymptotic probability of finding a tautology among all formulas.

This paper is a continuation of research concerning the *density of truth* in different logics. Until now, we were concentrated mostly on classical and intuitionistic logics. Especially, its  $\{\rightarrow\}$  and  $\{\rightarrow, \neg\}$  fragments with one variable were investigated (see [5], [9]). It is well known fact that implicational fragments of one variable of intuitionistic and classical logics are the same. Moreover, it is easy to observe that implicational-negational and monadic fragments of every intermediate logic is identical with that same fragment of intuitionistic one. It is already known (see [3]) that the *density* of that logic (and each intermediate as well) is more than 39%. We also know the *density* of implicational-negational and monadic fragment of classical logic. It is about 42%. That gives a very

large amount of intuitionistic tautologies among classical ones in the language consisting of signs of  $\{\rightarrow, \neg\}$  and one propositional variable.

It is natural to wish to investigate some important intermediate logics with respect to their densities. We will concentrate our attention on two of them: the Medvedev logic and Dummett's one. To distinguish them we will take their monadic fragments over richer language consisting of operators  $\{\rightarrow, \vee, \neg\}$ . Such fragments have models obtained by a simple modification of the Rieger-Nishimura lattice.

## 2 Intermediate logics

We consider the set of all formulas  $F$  built up from one variable  $p$  by means of operations  $\{\rightarrow, \vee, \neg\}$ . Our starting point is Medvedev's logic  $ML$  of finite problems. As it is known it is not finitely axiomatizable and might be characterized with help of Kreisel and Putnam's logic [4]. Recall that the logic  $KP$  of Kreisel and Putnam is the least intermediate logic with

$$(\neg\alpha \rightarrow (\beta \vee \gamma)) \rightarrow ((\neg\alpha \rightarrow \beta) \vee (\neg\alpha \rightarrow \gamma)) \quad (1)$$

Let  $F^{\{\neg\}}$  be the set of formulas defined as follows:

$$\neg\alpha \in F^{\{\neg\}} \Leftrightarrow \alpha \in F; \quad (2)$$

$$\alpha, \beta \in F^{\{\neg\}} \Rightarrow \alpha \rightarrow \beta, \alpha \vee \beta \in F^{\{\neg\}}. \quad (3)$$

The characterization is the following:

$$\alpha \in ML \Leftrightarrow (e(\alpha) \in KP, \text{ for every substitution } e : F \rightarrow F^{\{\neg\}})$$

Clearly  $KP \subset ML$ . The following formula

$$((\neg\neg\alpha \rightarrow \alpha) \rightarrow (\alpha \vee \neg\alpha)) \rightarrow (\neg\neg\alpha \vee \neg\alpha) \quad (4)$$

called Scott's law, belongs to  $ML$  and does not belong to  $KP$ , see [8].

Hence, the Lindenbaum algebra for the monadic fragment of  $ML$  is obtained by dividing the Rieger-Nishimura lattice by the filter generated just by the Scott's law. We will denote this algebra by  $\mathcal{M}$ . It is consisted of 9 following equivalence classes:

$$\begin{aligned} A &= [\neg(p \rightarrow p)]_{\equiv}, \\ B &= [p]_{\equiv}, \\ C &= [\neg p]_{\equiv}, \\ D &= [p \vee \neg p]_{\equiv}, \\ E &= [\neg\neg p]_{\equiv}, \\ G &= [\neg\neg p \vee (p \vee \neg p)]_{\equiv}, \\ H &= [\neg\neg p \rightarrow p]_{\equiv}, \\ J &= [(\neg\neg p \vee (p \vee \neg p)) \vee (\neg\neg p \rightarrow p)]_{\equiv}, \\ M &= [p \rightarrow p]_{\equiv}, \end{aligned}$$

and its diagram is as below:

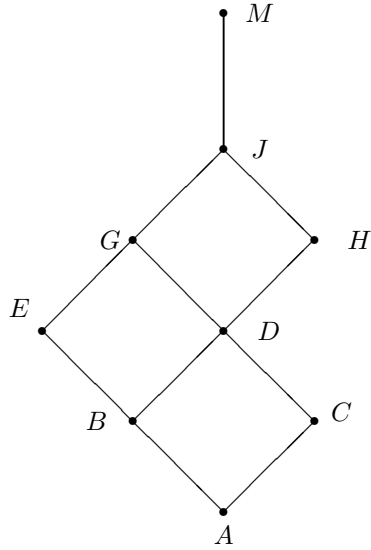


Figure 1.

### 3 Counting formulas and generating functions

In this section we set up the way of counting formulas with the established length.

way:

$$\begin{aligned}
 l(p) &= 1 \\
 l(\neg\phi) &= 1 + l(\phi) \\
 l(\phi \rightarrow \psi) &= l(\phi) + l(\psi) + 1 \\
 l(\phi \vee \psi) &= l(\phi) + l(\psi) + 1
 \end{aligned}$$

length  $n$ .

We will also consider some appropriate subclasses of  $F_n$ . For example if we have a class  $A \subset F$  then  $A_n = F_n \cap A$ .

class  $A_n$ .

**Lemma 1** *The number  $|F_n|$  of formulas from  $F_n$  is given by the recursion:*

$$|F_0| = 0, \quad |F_1| = 1, \quad (5)$$

$$|F_n| = |F_{n-1}| + 2 \sum_{i=1}^{n-2} |F_i| |F_{n-i-1}|. \quad (6)$$

*Proof.* Any formula of length  $n$  for  $n > 1$  is either a negation of some formula of length  $n-1$  for which the fragment  $|F_{n-1}|$  corresponds, or an implication (or disjunction) between some pairs of formulas of lengths  $i$  and  $n-i-1$ , respectively. Therefore the total number of such formulas is  $2 \sum_{i=1}^{n-2} |F_i| |F_{n-i-1}|$ .  $\square$

The main tool for dealing with asymptotics of sequences of numbers are *generating functions*, see for example [7].

Let  $A = (A_0, A_1, A_2, \dots)$  be a sequence of real numbers. It is in one-to-one correspondence to the formal power series  $\sum_{n=0}^{\infty} A_n z^n$ . Moreover, considering  $z$  as a complex variable, this series, as it is known from the theory of analytic functions, converges uniformly to a function  $f_A(z)$  in some open disc  $\{z \in \mathcal{C} : |z| < R\}$ , where  $R \geq 0$  is called its radius of convergence. So, with the sequence  $A$  we can associate a complex function  $f_A(z)$ , called the *ordinary generating function* for  $A$ , defined in a neighborhood of 0. This correspondence is one-to-one again (unless  $R = 0$ ), since the expansion of a complex function  $f(z)$ , analytic in a neighborhood of  $z_0$ , into a power series  $\sum_{n=0}^{\infty} A_n (z - z_0)^n$  is unique, and moreover, this series is the Taylor series, given by

$$A_n = \frac{1}{n!} \frac{d^n f}{dz^n}(z_0). \quad (7)$$

The above formula is a recursive one. To find a nonrecursive formula for  $A_n$  we take advantage from the following result due to Szegö [6] [Thm. 8.4], see also [7] [Thm. 5.3.2]. We need the following much simpler version of the Szegö lemma.

**Lemma 2** *Let  $v(z)$  be analytic in  $|z| < 1$  with  $z = 1$  being the only singularity at the circle  $|z| = 1$ . If  $v(z)$  in the vicinity of  $z = 1$  has an expansion of the form*

$$v(z) = \sum_{p \geq 0} v_p (1 - z)^{\frac{p}{2}}, \quad (8)$$

where  $p > 0$ , and the branch chosen above for the expansion equals  $v(0)$  for  $z = 0$ , then

$$[z^n]\{v(z)\} = v_1 \binom{1/2}{n} (-1)^n + O(n^{-2}). \quad (9)$$

The symbol  $[z^n]\{v(z)\}$  stands for the coefficient of  $z^n$  in the exponential series expansion of  $v(z)$ .

Now, we quote (without proof) some theorems which have appeared in [3]. They are the main tools for finding limits of the fraction  $\frac{a_n}{b_n}$  when generating functions for sequences  $a_n$  and  $b_n$  satisfy conditions of simplified Szegö Lemma 2.

**Lemma 3** Suppose two functions  $v(z)$  and  $w(z)$  satisfy assumptions of simplified Szegő theorem (Lemma 2) i.e. both  $v$  and  $w$  are analytic in  $|z| < 1$  with  $z = 1$  being the only singularity at the circle  $|z| = 1$ . Both  $v(z)$  and  $w(z)$  in the vicinity of  $z = 1$  have expansions of the form

$$v(z) = \sum_{p \geq 0} v_p (1-z)^{p/2},$$

$$w(z) = \sum_{p \geq 0} w_p (1-z)^{p/2},$$

then the limit of  $\frac{[z^n]\{v(z)\}}{[z^n]\{w(z)\}}$  exists and is given by the formula:

$$\lim_{n \rightarrow \infty} \frac{[z^n]\{v(z)\}}{[z^n]\{w(z)\}} = \frac{v_1}{w_1}$$

**Theorem 4** Suppose two functions  $v(z)$  and  $w(z)$  satisfy assumptions of simplified Szegő theorem (Lemma 2) i.e. both  $v$  and  $w$  are analytic in  $|z| < 1$  with  $z = 1$  being the only singularity at the circle  $|z| = 1$ . Both  $v(z)$  and  $w(z)$  in the vicinity of  $z = 1$  have expansions of the form

$$v(z) = \sum_{p \geq 0} v_p (1-z)^{p/2},$$

$$w(z) = \sum_{p \geq 0} w_p (1-z)^{p/2},$$

Suppose we have functions  $\tilde{v}$  and  $\tilde{w}$  satisfying  $\tilde{v}(\sqrt{1-z}) = v(z)$  and  $\tilde{w}(\sqrt{1-z}) = w(z)$  then the limit of  $\frac{[z^n]\{v(z)\}}{[z^n]\{w(z)\}}$  exists and is given by the formula:

$$\lim_{n \rightarrow \infty} \frac{[z^n]\{v(z)\}}{[z^n]\{w(z)\}} = \frac{(\tilde{v})'(0)}{(\tilde{w})'(0)} \quad (10)$$

## 4 Gluing of classes

In this section we do some preparations for determining the generating function for the class of tautologies of  $ML$ . First, we determine the generating function for the sequence of numbers  $|F_n|$ .

**Lemma 5** The generating function  $f$  for the numbers  $|F_n|$  is the following:

$$f(z) = \frac{1-z}{4} - \frac{\sqrt{z^2 - 10z + 1}}{4}. \quad (11)$$

*Proof.* The recurrence formula  $|F_n| = |F_{n-1}| + 2 \sum_{i=1}^{n-2} |F_i| |F_{n-i-1}|$  becomes the equality

$$f(z) = zf(z) + 2f^2(z) + z \quad (12)$$

since the recursion fragment  $\sum_{i=1}^{n-2} |F_i||F_{n-i-1}|$  corresponds exactly to multiplication of power series. The term  $|F_{n-1}|$  corresponds to the function  $zf(z)$ . The linear term  $z$  corresponds to the first non-zero coefficient in the power series of  $f$ . Solving the equation we get two possible solutions:  $f(z) = \frac{1-z}{4} - \frac{\sqrt{z^2-10z+1}}{4}$  or  $f(z) = \frac{1-z}{4} + \frac{\sqrt{z^2-10z+1}}{4}$ . We choose the first one, since it corresponds to the boundary condition  $f(0) = 0$  (see equation (5)).  $\square$

To find the generating functions for other classes of formulas it will be useful to have written the truth-table for the operations  $\{\rightarrow, \vee, \neg\}$  of the algebra  $\mathcal{M}$  presented in Figure 1.

$\rightarrow$	<i>A</i>	<i>B</i>	<i>C</i>	<i>D</i>	<i>E</i>	<i>G</i>	<i>H</i>	<i>J</i>	<i>M</i>	$\neg$
<i>A</i>	<i>M</i>	<i>M</i>	<i>M</i>	<i>M</i>	<i>M</i>	<i>M</i>	<i>M</i>	<i>M</i>	<i>M</i>	<i>M</i>
<i>B</i>	<i>C</i>	<i>M</i>	<i>C</i>	<i>M</i>	<i>M</i>	<i>M</i>	<i>M</i>	<i>M</i>	<i>M</i>	<i>C</i>
<i>C</i>	<i>E</i>	<i>E</i>	<i>M</i>	<i>M</i>	<i>E</i>	<i>M</i>	<i>M</i>	<i>M</i>	<i>M</i>	<i>E</i>
<i>D</i>	<i>A</i>	<i>E</i>	<i>C</i>	<i>M</i>	<i>E</i>	<i>M</i>	<i>M</i>	<i>M</i>	<i>M</i>	<i>A</i>
<i>E</i>	<i>C</i>	<i>H</i>	<i>C</i>	<i>H</i>	<i>M</i>	<i>M</i>	<i>H</i>	<i>M</i>	<i>M</i>	<i>C</i>
<i>G</i>	<i>A</i>	<i>B</i>	<i>C</i>	<i>H</i>	<i>E</i>	<i>M</i>	<i>H</i>	<i>M</i>	<i>M</i>	<i>A</i>
<i>H</i>	<i>A</i>	<i>E</i>	<i>C</i>	<i>G</i>	<i>E</i>	<i>G</i>	<i>M</i>	<i>M</i>	<i>M</i>	<i>A</i>
<i>J</i>	<i>A</i>	<i>B</i>	<i>C</i>	<i>D</i>	<i>E</i>	<i>G</i>	<i>H</i>	<i>M</i>	<i>M</i>	<i>A</i>
<i>M</i>	<i>A</i>	<i>B</i>	<i>C</i>	<i>D</i>	<i>E</i>	<i>G</i>	<i>H</i>	<i>J</i>	<i>M</i>	<i>A</i>

Table 1.

$\vee$	<i>A</i>	<i>B</i>	<i>C</i>	<i>D</i>	<i>E</i>	<i>G</i>	<i>H</i>	<i>J</i>	<i>M</i>
<i>A</i>	<i>A</i>	<i>B</i>	<i>C</i>	<i>D</i>	<i>E</i>	<i>G</i>	<i>H</i>	<i>J</i>	<i>M</i>
<i>B</i>	<i>B</i>	<i>B</i>	<i>D</i>	<i>D</i>	<i>E</i>	<i>G</i>	<i>H</i>	<i>J</i>	<i>M</i>
<i>C</i>	<i>C</i>	<i>D</i>	<i>C</i>	<i>D</i>	<i>G</i>	<i>G</i>	<i>H</i>	<i>J</i>	<i>M</i>
<i>D</i>	<i>D</i>	<i>D</i>	<i>D</i>	<i>D</i>	<i>G</i>	<i>G</i>	<i>H</i>	<i>J</i>	<i>M</i>
<i>E</i>	<i>E</i>	<i>E</i>	<i>G</i>	<i>G</i>	<i>E</i>	<i>G</i>	<i>J</i>	<i>J</i>	<i>M</i>
<i>G</i>	<i>G</i>	<i>G</i>	<i>G</i>	<i>G</i>	<i>G</i>	<i>G</i>	<i>J</i>	<i>J</i>	<i>M</i>
<i>H</i>	<i>H</i>	<i>H</i>	<i>H</i>	<i>H</i>	<i>J</i>	<i>J</i>	<i>H</i>	<i>J</i>	<i>M</i>
<i>J</i>	<i>J</i>	<i>J</i>	<i>J</i>	<i>J</i>	<i>J</i>	<i>J</i>	<i>J</i>	<i>J</i>	<i>M</i>
<i>M</i>	<i>M</i>	<i>M</i>	<i>M</i>	<i>M</i>	<i>M</i>	<i>M</i>	<i>M</i>	<i>M</i>	<i>M</i>

Table 2.

From these tables we can read the way of creating each tautology and any formula from any other class. For example the formulas from the class *A* are built up as implications between formulas from the classes *D, G, H, J, M* and *A*, negations of formulas from *D, G, H, J, M* and disjunctions of formulas from class *A*. But if we want to count the quantities of that formulas we would have to know the way of creating formulas from the other needed classes. This provides to obtaining a system of nine recurrent equations and consequently a system of nine non-linear equations with nine unknown generating functions. To avoid such complication we will consider step by step much simpler algebras obtained from the main quotient algebra  $\mathcal{M}$ .

First let us take the filter  $\{J, M\}$  and consider the new algebra  $\mathcal{M}/_{\{J, M\}}$ . Its diagram and the appropriate truth - tables are presented below.

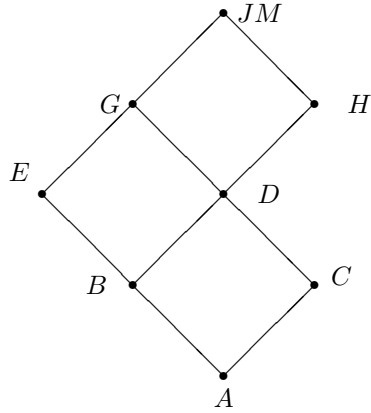


Figure 2

where

$$JM = J \cup M$$

We will repeat this abbreviation for other classes.

$\rightarrow$	<i>A</i>	<i>B</i>	<i>C</i>	<i>D</i>	<i>E</i>	<i>G</i>	<i>H</i>	<i>JM</i>	$\neg$
<i>A</i>	<i>JM</i>	<i>JM</i>	<i>JM</i>	<i>JM</i>	<i>JM</i>	<i>JM</i>	<i>JM</i>	<i>JM</i>	<i>JM</i>
<i>B</i>	<i>C</i>	<i>JM</i>	<i>C</i>	<i>JM</i>	<i>JM</i>	<i>JM</i>	<i>JM</i>	<i>JM</i>	<i>C</i>
<i>C</i>	<i>E</i>	<i>E</i>	<i>JM</i>	<i>JM</i>	<i>E</i>	<i>JM</i>	<i>JM</i>	<i>JM</i>	<i>E</i>
<i>D</i>	<i>A</i>	<i>E</i>	<i>C</i>	<i>JM</i>	<i>E</i>	<i>JM</i>	<i>JM</i>	<i>JM</i>	<i>A</i>
<i>E</i>	<i>C</i>	<i>H</i>	<i>C</i>	<i>H</i>	<i>JM</i>	<i>JM</i>	<i>H</i>	<i>JM</i>	<i>C</i>
<i>G</i>	<i>A</i>	<i>B</i>	<i>C</i>	<i>H</i>	<i>E</i>	<i>JM</i>	<i>H</i>	<i>JM</i>	<i>A</i>
<i>H</i>	<i>A</i>	<i>E</i>	<i>C</i>	<i>G</i>	<i>E</i>	<i>G</i>	<i>JM</i>	<i>JM</i>	<i>A</i>
<i>JM</i>	<i>A</i>	<i>B</i>	<i>C</i>	<i>D</i>	<i>E</i>	<i>G</i>	<i>H</i>	<i>JM</i>	<i>A</i>

Table 3.

$\vee$	<i>A</i>	<i>B</i>	<i>C</i>	<i>D</i>	<i>E</i>	<i>G</i>	<i>H</i>	<i>JM</i>
<i>A</i>	<i>A</i>	<i>B</i>	<i>C</i>	<i>D</i>	<i>E</i>	<i>G</i>	<i>H</i>	<i>JM</i>
<i>B</i>	<i>B</i>	<i>B</i>	<i>D</i>	<i>D</i>	<i>E</i>	<i>G</i>	<i>H</i>	<i>JM</i>
<i>C</i>	<i>C</i>	<i>D</i>	<i>C</i>	<i>D</i>	<i>G</i>	<i>G</i>	<i>H</i>	<i>JM</i>
<i>D</i>	<i>D</i>	<i>D</i>	<i>D</i>	<i>D</i>	<i>G</i>	<i>G</i>	<i>H</i>	<i>JM</i>
<i>E</i>	<i>E</i>	<i>E</i>	<i>G</i>	<i>G</i>	<i>E</i>	<i>G</i>	<i>JM</i>	<i>JM</i>
<i>G</i>	<i>G</i>	<i>G</i>	<i>G</i>	<i>G</i>	<i>G</i>	<i>G</i>	<i>JM</i>	<i>JM</i>
<i>H</i>	<i>H</i>	<i>H</i>	<i>H</i>	<i>H</i>	<i>JM</i>	<i>JM</i>	<i>H</i>	<i>JM</i>
<i>JM</i>	<i>JM</i>	<i>JM</i>	<i>JM</i>	<i>JM</i>	<i>JM</i>	<i>JM</i>	<i>JM</i>	<i>JM</i>

Table 4.

Now, we take the filter  $\{H, J, M\}$  and consider the new algebra  $\mathcal{M}/_{\{H, J, M\}}$ . Its diagram and the appropriate truth - tables are the following:

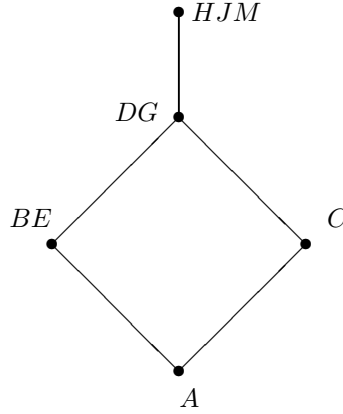


Figure 3.

$\rightarrow$	<i>A</i>	<i>BE</i>	<i>C</i>	<i>DG</i>	<i>HJM</i>	$\neg$
<i>A</i>	<i>HJM</i>	<i>HJM</i>	<i>HJM</i>	<i>HJM</i>	<i>HJM</i>	<i>HJM</i>
<i>BE</i>	<i>C</i>	<i>HJM</i>	<i>C</i>	<i>HJM</i>	<i>HJM</i>	<i>C</i>
<i>C</i>	<i>BE</i>	<i>BE</i>	<i>HJM</i>	<i>HJM</i>	<i>HJM</i>	<i>BE</i>
<i>DG</i>	<i>A</i>	<i>BE</i>	<i>C</i>	<i>HJM</i>	<i>HJM</i>	<i>A</i>
<i>HJM</i>	<i>A</i>	<i>BE</i>	<i>C</i>	<i>DG</i>	<i>HJM</i>	<i>A</i>

Table 5.

$\vee$	<i>A</i>	<i>BE</i>	<i>C</i>	<i>DG</i>	<i>HJM</i>
<i>A</i>	<i>A</i>	<i>BE</i>	<i>C</i>	<i>DG</i>	<i>HJM</i>
<i>BE</i>	<i>BE</i>	<i>BE</i>	<i>DG</i>	<i>DG</i>	<i>HJM</i>
<i>C</i>	<i>C</i>	<i>DG</i>	<i>C</i>	<i>DG</i>	<i>HJM</i>
<i>DG</i>	<i>DG</i>	<i>DG</i>	<i>DG</i>	<i>DG</i>	<i>HJM</i>
<i>HJM</i>	<i>HJM</i>	<i>HJM</i>	<i>HJM</i>	<i>HJM</i>	<i>HJM</i>

Table 6.

We also divide the algebra  $\mathcal{M}$  by the filter  $L = \{G, J, M\}$ :



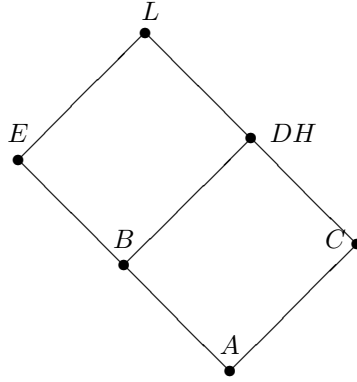


Figure 4.

The truth tables are the following:

$\rightarrow$	$A$	$B$	$C$	$DH$	$E$	$L$	$\neg$
$A$	$L$	$L$	$L$	$L$	$L$	$L$	$L$
$B$	$C$	$L$	$C$	$L$	$L$	$L$	$C$
$C$	$E$	$E$	$L$	$L$	$E$	$L$	$E$
$DH$	$A$	$E$	$C$	$L$	$E$	$L$	$A$
$E$	$C$	$DH$	$C$	$DH$	$L$	$L$	$C$
$L$	$A$	$B$	$C$	$DH$	$E$	$L$	$A$

Table 7.

$\vee$	$A$	$B$	$C$	$DH$	$E$	$L$
$A$	$A$	$B$	$C$	$DH$	$E$	$L$
$B$	$B$	$B$	$DH$	$DH$	$E$	$L$
$C$	$C$	$DH$	$C$	$DH$	$L$	$L$
$DH$	$DH$	$DH$	$DH$	$DH$	$L$	$L$
$E$	$E$	$E$	$L$	$L$	$E$	$L$
$L$	$L$	$L$	$L$	$L$	$L$	$L$

Table 8.

**Lemma 6** *The matrix described in Tables 7 and 8 is a matrix for the linear Dummett's logic of implication, disjunction and negation with one variable*<sup>1</sup>.

*Proof.* Proof of that fact is poorly semantical, analogous to the one of Lemma 7 in [3].

The next divisions are the following - we divide the algebra  $\mathcal{M}$  by the filters  $\{E, G, J, M\}$  and  $K = \{D, G, H, J, M\}$  and obtain:

<sup>1</sup>Linear calculus  $LC$  was studied in [1] by Dummett. Syntactically it is obtained by adding the axiom  $(p \rightarrow q) \vee (q \rightarrow p)$  to axioms of intuitionistic logic.



Figure 5.

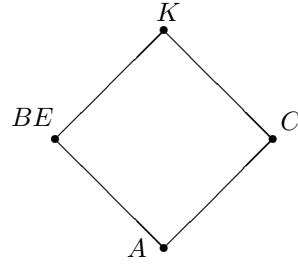


Figure 6.

The operations  $\{\rightarrow, \vee, \neg\}$  in the new algebras are given by the following truth-tables:

$\rightarrow$	<i>AC</i>	<i>BDH</i>	<i>EGJM</i>	$\neg$
<i>AC</i>	<i>EGJM</i>	<i>EGJM</i>	<i>EGJM</i>	<i>EGJM</i>
<i>BDH</i>	<i>AC</i>	<i>EGJM</i>	<i>EGJM</i>	<i>AC</i>
<i>EGJM</i>	<i>AC</i>	<i>BDH</i>	<i>EGJM</i>	<i>AC</i>

Table 9.

$\vee$	<i>AC</i>	<i>BDH</i>	<i>EGJM</i>
<i>AC</i>	<i>AC</i>	<i>BDH</i>	<i>EGJM</i>
<i>BDH</i>	<i>BDH</i>	<i>BDH</i>	<i>EGJM</i>
<i>EGJM</i>	<i>EGJM</i>	<i>EGJM</i>	<i>EGJM</i>

Table 10.

$\rightarrow$	$A$	$BE$	$C$	$K$	$\neg$
$A$	$K$	$K$	$K$	$K$	$K$
$BE$	$C$	$K$	$C$	$K$	$C$
$C$	$BE$	$BE$	$K$	$K$	$BE$
$K$	$A$	$BE$	$C$	$K$	$A$

Table 11.

$\vee$	$A$	$BE$	$C$	$K$
$A$	$A$	$BE$	$C$	$K$
$BE$	$BE$	$BE$	$K$	$K$
$C$	$C$	$K$	$C$	$K$
$K$	$K$	$K$	$K$	$K$

Table 12.

As we can observe, the first truth table describes operations in the Gödel 3 valued matrix, while the second one is a matrix of all valuations associated with the standard classical logic of one variable.

**Lemma 7** *The matrix described in Tables 11 and 12 is a matrix for the classical logic of implication, disjunction and negation with one variable.*

*Proof.* Proof of that fact is analogous to the one of Lemma 12 in [3].  $\square$

The last two divisions are the following: we take two filters  $N = \{C, D, G, H, J, M\}$  and  $P = \{B, D, E, G, H, J, M\}$ . The new quotient algebras  $\mathcal{M}/_N$  and  $\mathcal{M}/_P$  have the following diagrams:

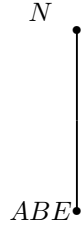


Figure 7.

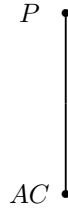


Figure 8.

The operations  $\{\rightarrow, \vee, \neg\}$  are characterized by the following truth-tables:

$\rightarrow$	$ABE$	$N$	$\neg$
$ABE$	$N$	$N$	$N$
$N$	$ABE$	$N$	$ABE$

Table 13.

$\rightarrow$	$ABE$	$N$
$ABE$	$ABE$	$N$
$N$	$N$	$N$

Table 14.

$\rightarrow$	$AC$	$P$	$\neg$
$AC$	$P$	$P$	$P$
$P$	$AC$	$P$	$AC$

Table 15.

$\rightarrow$	$AC$	$P$
$AC$	$AC$	$P$
$P$	$P$	$P$

Table 16.

## 5 Calculating generating functions

Now, we are ready to deal with generating functions. We start with analyzing the algebra  $\mathcal{M}/N$ .

**Lemma 8** *The numbers  $|ABE_n|$  are given by the following mutual recursions:*

$$|ABE_0| = 0, \quad |ABE_1| = 1, \quad (13)$$

$$|ABE_n| = \sum_{i=1}^{n-2} |F_i| |ABE_{n-i-1}| + |N_{n-1}|, \quad (14)$$

*Proof.* From Tables 13 and 14 we see that formulas from class  $ABE$  can be obtained as implications of formulas from classes  $N$  and  $ABE$  - this gives us the recurrence  $\sum_{i=1}^{n-2} |N_i| |ABE_{n-i-1}|$ , disjunctions of formulas from class  $ABE$  - hence we have  $\sum_{i=1}^{n-2} |ABE_i| |ABE_{n-i-1}|$ , or negations of ones from the class  $N$ , which gives the part  $|N_{n-1}|$ . Hence we have:

$$|ABE_n| = \sum_{i=1}^{n-2} |N_i| |ABE_{n-i-1}| + \sum_{i=1}^{n-2} |ABE_i| |ABE_{n-i-1}| + |N_{n-1}|.$$

From disjointness of considered classes we have  $|ABE_i| + |N_i| = |F_i|$  and hence we have (14). Because  $p \in ABE$  then we have (13).  $\square$

**Lemma 9** *The generating function  $f_{ABE}$  for sequence of numbers  $|ABE_n|$  is:*

$$f_{ABE}(z) = \frac{zf(z) + z}{1 - f(z) + z} \quad (15)$$

$$\text{where } f(z) = \frac{1-z}{4} - \frac{\sqrt{z^2 - 10z + 1}}{4}$$

*Proof.* The part  $\sum_{i=1}^{n-2} |F_i| |ABE_{n-i-1}|$  of the recurrence (14) corresponds to multiplication of power series and then gives multiplication of generating functions:  $f(z)f_{ABE}(z)$ . The term  $|N_{n-1}| = |F_i| - |ABE_i|$  corresponds to the function  $z(f(z) - f_{ABE}(z))$ . Hence the recurrence (14) gives the following equality between the appropriate generating functions:

$$f_{ABE}(z) = f(z)f_{ABE}(z) + z(f(z) - f_{ABE}(z)) + z \quad (16)$$

The linear term  $z$  in (16) corresponds to the first non-zero coefficient in the power series of  $f_{ABE}$ . A basic transformation gives us (15).  $\square$

Analogously we determine the appropriate recurrence and generating function for the numbers  $|AC_n|$ .

**Lemma 10** *The numbers  $|AC_n|$  are given by the following mutual recursions:*

$$|AC_0| = |AC_1| = 0, \quad |AC_2| = 1, \quad (17)$$

$$|AC_n| = \sum_{i=1}^{n-2} |F_i| |AC_{n-i-1}| + |P_{n-1}|, \quad (18)$$

*Proof.* The above recurrence follows from Tables 15 and 16 analogously to the one in Lemma 8.  $\square$

**Lemma 11** *The generating function  $f_{AC}$  for sequence of numbers  $|AC_n|$  is the following:*

$$f_{AC}(z) = \frac{zf(z)}{1-f(z)+z} \quad (19)$$

$$\text{where } f(z) = \frac{1-z}{4} - \frac{\sqrt{z^2-10z+1}}{4}$$

*Proof.* The recurrence (18) can be translated into following equation:

$$f_{AC}(z) = f(z)f_{AC}(z) + z(f(z) - f_{AC}(z)) \quad (20)$$

which gives us (19). There is no linear term in (20) because  $p \notin AC$ .  $\square$

In the same manner we can determine generating functions connected with other classes of formulas. For simplicity of notation we will omit the argument  $(z)$  in all the functions which will arrive hereafter. Now, we determine the generating function for the class of classical tautologies  $K$ .

**Lemma 12** *The generating function  $f_K$  for the numbers  $|K_n|$  is the following:*

$$f_K = \frac{1}{16} \left( 4 - 4z + \frac{12f(-1+f+z) + 12z}{-1+f-z} - S \right), \quad (21)$$

*where*

$$f(z) = \frac{1-z}{4} - \frac{\sqrt{z^2-10z+1}}{4}$$

$$S = \frac{\sqrt{2}\sqrt{8-z+20z^2+31z^3+8z^4} + f(-15-31z-z^2+15z^3)}{1-f+z}.$$

*Proof.* From Tables 11 and 12 we have the following recurrence for the numbers  $|A_n|$  hold:

$$|A_0| = |A_1| = |A_2| = |A_3| = 0, \quad |A_4| = 1, \quad (22)$$

$$|A_n| = \sum_{i=1}^{n-2} |K_i| |A_{n-i-1}| + \sum_{i=1}^{n-2} |A_i| |A_{n-i-1}| + |K_{n-1}|, \quad (23)$$

This recurrence can be translated into equation:

$$f_A = f_K f_A + f_A^2 + z f_K \quad (24)$$

with two unknown functions  $f_A$  and  $f_K$ .

From the other side, from definitions of classes  $AC$ ,  $ABE$  and  $K$  we have:

$$f_A = f_{AC} + f_K + f_{ABE} - f \quad (25)$$

From the system of equations:

$$\begin{cases} (24) \\ (25) \end{cases}$$

we obtain the quadratic equation:

$$2f_K^2 + 3f_K(f_{AC} + f_{ABE} - f) + (f_{AC} + f_{ABE} - f)^2 - f_{AC} - f_{ABE} + f = 0$$

After solving it with the boundary condition  $f_K(0) = 0$  and taking into consideration the equalities (19), (15), (11) and intensive simplification we get (21).  $\square$

The next important generating function is the one of class of linear tautologies  $L$  (see Figure 4). To determine it we must first consider Figure 5 and determine the generating function for the class of formulas  $BDH$ .

**Lemma 13** *The generating function  $f_{BDH}$  for the sequence of numbers  $|BDH_n|$  is:*

$$f_{BDH} = \frac{2z(1-f+z)}{-3f-5zf+z+2}, \quad (26)$$

where

$$f(z) = \frac{1-z}{4} - \frac{\sqrt{z^2-10z+1}}{4}$$

*Proof.* It follows easily from Tables 9 and 10 that the numbers  $|BDH_n|$  are given by the following recursion:

$$\begin{aligned} |BDH_0| &= 0, \quad |BDH_1| = 1, \\ |BDH_n| &= \sum_{i=1}^{n-2} |EGJM_i| |BDH_{n-i-1}| + 2 \sum_{i=1}^{n-2} |BDH_i| |AC_{n-i-1}| + \\ &+ \sum_{i=1}^{n-2} |BDH_i| |BDH_{n-i-1}|. \end{aligned} \quad (27)$$

The above recurrence leads to the equality:

$$f_{BDH} = f_{EGJM} f_{BDH} + 2f_{BDH} f_{AC} + f_{BDH}^2 + z \quad (28)$$

From disjointness of the appropriate classes we have that:

$$f_{EGJM} = f - f_{BDH} - f_{AC}. \quad (29)$$

and after application (29) to the equation (28) we obtain as follows:

$$f_{BDH} = \frac{z}{1-f_{AC}-f} \quad (30)$$

which after a suitable simplification gives us (26).  $\square$

The next needed function is the function  $f_B$ . We take into consideration Figure 4.

**Lemma 14** *The generating function  $f_B$  for the sequence of numbers  $|B_n|$  is the following:*

$$f_B = \frac{12 + S + 4z + U - 16X - 16\sqrt{-8z + (-1 - 2f + Y + X)^2}}{64}, \quad (31)$$

where

$$\begin{aligned} f(z) &= \frac{1-z}{4} - \frac{\sqrt{z^2 - 10z + 1}}{4} \\ S &= \frac{\sqrt{2}\sqrt{8-z+20z^2+31z^3+8z^4} + f(-15-31z-z^2+15z^3)}{1-f+z} \\ U &= \frac{-8z - f(12f-2S) + 2(4-4z^2-S(1+z))}{-1+f-z} \\ X &= \frac{2z(1-f+z)}{-2-z+f(3+5z)} \\ Y &= \frac{-14z - 58fz + 3(-4+S-f(2+S)+Sz+4z^2)}{16(-1+f-z)} \end{aligned} \quad (32)$$

*Proof.* From Tables 7 and 8 we get the following recurrence concerning the numbers  $|B_n|$ :

$$\begin{aligned} |B_0| &= 0, \quad |B_1| = 1, \\ |B_n| &= \sum_{i=1}^{n-2} |L_i| |B_{n-i-1}| + 2 \sum_{i=1}^{n-2} |B_i| |A_{n-i-1}| + \\ &\quad + \sum_{i=1}^{n-2} |B_i| |B_{n-i-1}|. \end{aligned} \quad (33)$$

The above recurrence leads to the equality:

$$f_B = f_L f_B + 2f_B f_A + f_B^2 + z. \quad (34)$$

We also have:

$$\begin{aligned} f_L &= f_{EGJM} - f_E, \quad f_{EGJM} = f - f_{BDH} - f_{AC} \quad f_E = f_{BE} - f_B, \\ f_{BE} &= f - f_{AC} - f_K, \quad f_A = f_{AC} + f_K + f_{ABE} - f. \end{aligned}$$

After a suitable substitution and simplification we get:

$$f_L = f_K + f_B - f_{BDH}. \quad (35)$$

We apply (35) to (34) and after simplification obtain the following quadratic equation:

$$2f_B^2 + f_B(2f_{AC} + 3f_K + 2f_{ABE} - 2f - f_{BDH} - 1) + z = 0. \quad (36)$$

Solving (36) with the boundary condition  $f_B(0) = 0$  and a suitable substitution of (19), (21), (15) and (26) give us (31).  $\square$

Now, we are ready to determine the generating function of linear tautologies. Let us take advantage of Lemma 14 and (35).

**Corollary 15** *The generating function  $f_L$  for the sequence of numbers  $|L_n|$  is the following:*

$$f_L = \frac{1}{64} \left( 28 - 12z - 3S + U + 48X + 48 \frac{f(-1+f+z) + z}{-1+f-z} - 16\sqrt{-8z + (-1 - 2f + Y + X)^2} \right), \quad (37)$$

where

$$f = \frac{1-z}{4} - \frac{\sqrt{z^2 - 10z + 1}}{4},$$

$$S = \frac{\sqrt{2}\sqrt{8-z+20z^2+31z^3+8z^4} + f(-15-31z-z^2+15z^3)}{1-f+z}, \quad (38)$$

$$U = \frac{-8z - f(12f - 2S) + 2(4 - 4z^2 - S(1+z))}{-1+f-z}, \quad (39)$$

$$X = \frac{2z(1-f+z)}{-2-z+f(3+5z)}, \quad (40)$$

$$Y = \frac{-14z - 58fz + 3(-4 + S - f(2+S) + Sz + 4z^2)}{16(-1+f-z)}. \quad (41)$$

The next step will concern the lattice presented in Figure 3. We notice that:

**Lemma 16** *The generating function  $f_{DG}$  for the numbers  $|DG_n|$  is the following:*

$$f_{DG} = \frac{T}{8(-1+f-z)} * \frac{(T+16z)}{-20+S(1+z-f)+z(6+4z)+f(30+50z)} \quad (42)$$

where

$$f = \frac{1-z}{4} - \frac{\sqrt{z^2 - 10z + 1}}{4}$$

$$S = \frac{\sqrt{2}\sqrt{8-z+20z^2+31z^3+8z^4} + f(-15-31z-z^2+15z^3)}{1-f+z}$$

$$T = -4 + S - 2z + Sz + 4z^2 + f(6 - S + 10z)$$

*Proof.* From Tables 5 and 6 we get the following recurrence concerning the numbers  $|DG_n|$ :

$$|DG_0| = |DG_1| = |DG_2| = |DG_3| = 0, \quad |DG_4| = 1,$$

$$|DG_n| = \sum_{i=1}^{n-2} |HJM_i| |DG_{n-i-1}| + 2 \sum_{i=1}^{n-2} (|A_i| + |BE_i| + |C_i|) |DG_{n-i-1}| +$$



$$+ \sum_{i=1}^{n-2} |DG_i| |DG_{n-i-1}| + 2 \sum_{i=1}^{n-2} |C_i| |BE_{n-i-1}|. \quad (43)$$

The above recurrence gives the equality:

$$f_{DG} = f_{HJM} f_{DG} + 2(f_A + f_{BE} + f_C) f_{DG} + f_{DG}^2 + 2f_C f_{BE}. \quad (44)$$

Because  $f_{HJM} = f_K - f_{DG}$ ,  $f_A + f_{BE} + f_C = f - f_K$ ,  $f_C = f - f_K - f_{ABE}$  and  $f_{BE} = f - f_{AC} - f_K$  then we obtain the following linear equation with respect the function  $f_{DG}$ :

$$f_{DG}(2f - f_K - 1) + 2(f - f_K - f_{ABE})(f - f_{AC} - f_K) = 0. \quad (45)$$

This equation after suitable substitution and simplification give us (42).  $\square$

Now, we are very close to achieve the main goal of that paper. All that is left for us to do, is to determine two generating functions - one connected with Figure 2 (we will choose  $f_D$ ) and the last one which is the generating function for the class of tautologies of Medvedev's calculus. Let us consider Figure 2.

**Lemma 17** *The generating function  $f_D$  for the numbers  $|D_n|$  is the following:*

$$f_D = \frac{1}{4} (1 - 2f_{AC} - f_L + f_{DG} - 2f_B - \sqrt{16(f_{ABE} - f + f_K)f_B + (-1 + 2f_{AC} + f_L - f_{DG} + 2f_B)^2}) \quad (46)$$

where functions  $f$ ,  $f_{ABE}$ ,  $f_{AC}$ ,  $f_K$ ,  $f_B$ ,  $f_L$ ,  $f_{DG}$  are defined by (11), (15), (19), (21), (31), (37) and (42).

*Proof.* Tables 3 and 4 give us the following recurrence:

$$\begin{aligned} |D_0| &= |D_1| = |D_2| = |D_3| = 0, \quad |D_4| = 1, \\ |D_n| &= 2 \left( \sum_{i=1}^{n-2} |B_i| |C_{n-i-1}| + \sum_{i=1}^{n-2} |D_i| (|A_{n-i-1}| + |C_{n-i-1}| + |B_{n-i-1}|) \right) + \\ &\quad \sum_{i=1}^{n-2} |JM_i| |D_{n-i-1}| + \sum_{i=1}^{n-2} |D_i| |D_{n-i-1}|. \end{aligned} \quad (47)$$

This recurrence can be translated as follows:

$$f_D = f_{JM} f_D + 2(f_B f_C + (f_A + f_C + f_B) f_D) + f_D^2. \quad (48)$$

We know also that  $f_{JM} = f_L - f_{DG} + f_D$ ,  $f_C = f - f_K - f_{ABE}$  and  $f_A + f_C + f_B = f_{AC} + f_B$ .

Therefore we obtain the following quadratic equation with respect to the function  $f_D$ :

$$2f_D^2 + f_D(f_L - f_{DG} + 2(f_{AC} + f_B) - 1) + 2f_B(f - f_K - f_{ABE}) = 0. \quad (49)$$

After solving it with the condition  $f_D(0) = 0$  we get (46). For simplicity we have expressed  $f_D$  in terms of before determined generating functions.  $\square$

Now, we are prepare to determine the generating function characterizing tautologies of Medvedev's logic. To do that we choose the class  $J_n$  because it is the one which distinguishes the algebras  $\mathcal{M}$  and  $\mathcal{M}/\{JM\}$ .

**Lemma 18** *The generating function  $f_J$  for the numbers  $|J_n|$  is the following:*

$$f_J = \frac{2(f - f_{BDH} - f_{AC} - f_L - f_D + f_{DG})(f_K - f_L - f_D)}{1 + f_D - 2f + f_L - f_{DG}} \quad (50)$$

where functions  $f, f_{AC}, f_K, f_{BDH}, f_L, f_{DG}, f_D$  are defined by (11), (19), (21), (26), (37), (42) and (46).

*Proof.* Tables 1 and 2 give us the following recurrence:

$$\begin{aligned} |J_0| &= \dots = |J_8| = 0, \quad |J_9| = 1, \\ |J_n| &= \sum_{i=1}^{n-2} |M_i| |J_{n-i-1}| + 2 \left( \sum_{i=1}^{n-2} |H_i| (|E_{n-i-1}| + |G_{n-i-1}|) + \right. \\ &\quad \left. \sum_{i=1}^{n-2} |J_i| (|F_{n-i-1}| - |J_{n-i-1}| - |M_{n-i-1}|) \right) + \sum_{i=1}^{n-2} |J_i| |J_{n-i-1}|. \end{aligned} \quad (51)$$

This recurrence can be translated as follows:

$$f_J = f_M f_J + 2(f_H(f_E + f_G) + f_J(f - f_J - f_M)) + f_J^2. \quad (52)$$

After taking advantage of the following equalities  $f_M = f_{JM} - f_J$ ,  $f_{JM} = f_L + f_D - f_{DG}$ ,  $f_E + f_G = f - f_{BDH} - f_{AC} - f_{JM}$   $f_H = f_K - f_L - f_D$  we obtain a linear equation which after solving gives us (50).  $\square$

**Corollary 19** *The generating function  $f_M$  for the sequence of numbers  $|M_n|$  is the following:*

$$f_M = f_L + f_D - f_{DG} - f_J \quad (53)$$

where functions  $f_L, f_{DG}, f_D$  and  $f_J$  are defined by (37), (42), (46) and (50).

## 6 Counting asymptotic densities

In this section we do some calculations concerning singularities of the investigated generating functions. First, let us observe that:

**Lemma 20**  $z_0 = 5 - 2\sqrt{6}$  is the only singularity of  $f, f_K, f_L$  and  $f_M$  located in  $|z| \leq 5 - 2\sqrt{6}$ .

*Proof.* It is easy to observe the function  $f(z)$  has only singularities at  $z = 5 - 2\sqrt{6}$  and  $z = 5 + 2\sqrt{6}$ . To make sure the functions  $f_K$ ,  $f_L$  and  $f_M$  have the nearest one at  $z = 5 - 2\sqrt{6}$ , we had to solve the following complicated equations:

$$\begin{aligned}
-1 + f - z &= 0 \\
S &= 0 \\
-8z + (-1 - 2f + Y + X)^2 &= 0 \\
-20 + S(1 + z - f) + z(6 + 4z + f(30 + 50z)) &= 0 \\
16(f_{ABE} - f + f_K)f_B + (-1 + 2f_{AC} + f_L - f_{DG} + 2f_B)^2 &= 0 \\
-3f - 5zf + z + 2 &= 0 \\
1 + f_D - 2f + f_L - f_{DG} &= 0
\end{aligned}$$

where functions  $f$ ,  $f_{AC}$ ,  $f_{ABE}$ ,  $f_K$ ,  $f_{DG}$ ,  $f_B$ ,  $f_L$ ,  $f_D$ ,  $X$ ,  $Y$ ,  $S$  are defined by (11), (19), (15), (21), (42), (31), (37), (46), (40), (41), (38).

To do that we had to use extensively *the Mathematica* package and it occurred that all solutions which are different from  $z = 5 - 2\sqrt{6}$  are situated outside the disc  $|z| \leq 5 - 2\sqrt{6}$ .  $\square$

To apply the Szegő Lemma we have to have functions which are analytic in the open disc  $|z| < 1$ , and the nearest singularity is at  $z_0 = 1$ . For that purpose we are going to calibrate functions  $f$ ,  $f_K$ ,  $f_L$  and  $f_M$  in the following way:

$$\begin{aligned}
\widehat{f}(z) &= f\left(\frac{z}{5-2\sqrt{6}}\right) & \widehat{f}_K(z) &= f_K\left(\frac{z}{5-2\sqrt{6}}\right) \\
\widehat{f}_L(z) &= f_L\left(\frac{z}{5-2\sqrt{6}}\right) & \widehat{f}_M(z) &= f_M\left(\frac{z}{5-2\sqrt{6}}\right).
\end{aligned}$$

It is not essential for our task to express the above functions in explicit forms. We only note that the relations between power series of the appropriate functions are such as  $[z^n]\{f(z)\} = ([z^n]\{\widehat{f}(z)\})(5 - 2\sqrt{6})^n$ .

**Corollary 21**  $z_0 = 1$  is the only singularity of  $\widehat{f}$ ,  $\widehat{f}_K$ ,  $\widehat{f}_L$  and  $\widehat{f}_M$  located in  $|z| \leq 1$ .

**Theorem 22** Expansions of functions  $\widehat{f}$  and  $\widehat{f}_K$  in a neighborhood of  $z = 1$  are as follows:

$$\begin{aligned}
\widehat{f}(z) &= f_0 + f_1\sqrt{1-z} + \dots \\
\widehat{f}_K(z) &= k_0 + k_1\sqrt{1-z} + \dots
\end{aligned}$$

where

$$f_0 = \frac{-4 + 2\sqrt{6}}{4}, \quad f_1 = -\frac{1}{4}\sqrt{-48 + 20\sqrt{6}}, \quad k_0 \approx 0.06468\dots, \quad k_1 \approx -0.16307\dots$$

*Proof.* The above coefficients were found using *the Mathematica* package.  $\square$   
Analogously we count the appropriate coefficients of generating functions of linear tautologies and Medvedev's ones.

**Theorem 23** *Expansions of functions  $\widehat{f}_L$  and  $\widehat{f}_M$  in a neighborhood of  $z = 1$  are as follows:*

$$\begin{aligned}\widehat{f}_L(z) &= l_0 + l_1\sqrt{1-z} + \dots \\ \widehat{f}_M(z) &= m_0 + m_1\sqrt{1-z} + \dots\end{aligned}$$

where

$$l_0 \approx 0.05534\dots, \quad l_1 = -0.14583\dots, \quad m_0 \approx 0.054511\dots, \quad m_1 \approx -0.14279\dots$$

Now, we can calculate the density of implicational-disjunctive-negational parts of classical, linear and Medvedev's logic of one variable. By application of the Szegő lemma, Lemma 3 and Theorem 4 we get as follows:

**Theorem 24**

$$\begin{aligned}\mu(K) &= \lim_{n \rightarrow \infty} \frac{|K_n|}{|F_n|} = \lim_{n \rightarrow \infty} \frac{(k_1 \binom{1/2}{n} (-1)^n + O(n^{-2}))(5 - 2\sqrt{6})^n}{(f_1 \binom{1/2}{n} (-1)^n + O(n^{-2}))(5 - 2\sqrt{6})^n} \\ &= \lim_{n \rightarrow \infty} \frac{k_1}{f_1} (1 + o(1)) = \frac{k_1}{f_1} \approx 65.56\%\end{aligned}$$

**Theorem 25**

$$\begin{aligned}\mu(L) &= \lim_{n \rightarrow \infty} \frac{|L_n|}{|F_n|} = \lim_{n \rightarrow \infty} \frac{(l_1 \binom{1/2}{n} (-1)^n + O(n^{-2}))(5 - 2\sqrt{6})^n}{(f_1 \binom{1/2}{n} (-1)^n + O(n^{-2}))(5 - 2\sqrt{6})^n} \\ &= \lim_{n \rightarrow \infty} \frac{l_1}{f_1} (1 + o(1)) = \frac{l_1}{f_1} \approx 58.63\%\end{aligned}$$

**Theorem 26**

$$\begin{aligned}\mu(M) &= \lim_{n \rightarrow \infty} \frac{|M_n|}{|F_n|} = \lim_{n \rightarrow \infty} \frac{(m_1 \binom{1/2}{n} (-1)^n + O(n^{-2}))(5 - 2\sqrt{6})^n}{(f_1 \binom{1/2}{n} (-1)^n + O(n^{-2}))(5 - 2\sqrt{6})^n} \\ &= \lim_{n \rightarrow \infty} \frac{m_1}{f_1} (1 + o(1)) = \frac{m_1}{f_1} \approx 57.41\%\end{aligned}$$

The results presented above can be compared with analogous one concerning the density of implicational-negational parts of linear and classical logics with one variable. As it was said in **Introduction** they amount to 39% and 42%. Then one can notice the operation of disjunction is in fact a truth 'makers'. From the other side we can see that the quantitative difference between considered fragments of Medvedev's and linear calculi are almost imperceptibly.

**Theorem 27** *The relative probability of finding a tautology of Medvedev's logic among linear ones is about 98 %.*

*Proof.* From the known asymptotics  $\lim_{n \rightarrow \infty} \frac{|M_n|}{|F_n|}$  and  $\lim_{n \rightarrow \infty} \frac{|L_n|}{|F_n|}$  we get

$$\lim_{n \rightarrow \infty} \frac{|M_n|}{|L_n|} = \frac{\lim_{n \rightarrow \infty} \frac{|M_n|}{|F_n|}}{\lim_{n \rightarrow \infty} \frac{|L_n|}{|F_n|}} = \frac{0.5741 \dots}{0.5863 \dots} \approx 98\%.$$

□

Finally, the above results can be employed to calculate the size of fragment of Dummett's logics inside classical one.

**Theorem 28** *The relative probability of finding a linear tautology among classical ones is more than 89 %.*

*Proof.* We already know asymptotics  $\lim_{n \rightarrow \infty} \frac{|L_n|}{|F_n|}$  and  $\lim_{n \rightarrow \infty} \frac{|K_n|}{|F_n|}$  therefore

$$\lim_{n \rightarrow \infty} \frac{|L_n|}{|K_n|} = \frac{\lim_{n \rightarrow \infty} \frac{|L_n|}{|F_n|}}{\lim_{n \rightarrow \infty} \frac{|K_n|}{|F_n|}} = \frac{0.5863 \dots}{0.6556 \dots} \approx 89\%.$$

□

Let us compare the last result with the one concerning the  $\{\rightarrow, \neg\}$  fragment of the monadic linear and classical logics. In [2] it is shown the relative probability of finding a linear tautology among classical ones of such language is more than 93%. This can be commented that the operations of disjunctions (linear and classical) play an important role in distinction between the considered logics.

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